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Fakultät Informatik Institut Software- und Multimediatechnik, Lehrstuhl Softwaretechnologie

Life-by-Wire

with Cyber-Physical Systems

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Technische Universität Dresden
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Linköping, May 2, 2012



ResUbic

ResUbic Lab
<http://www.resubic.org>

- Chap.1 The presence: Fly-by-wire
- Chap.2 The near future: Help-by-wire
 - Protect-by-wire
 - Rescue-by-wire
 - Serve-by-wire
- Chap.3 The mid-term future: Move-by-wire
 - Drive-by-wire
 - Move-by-wire
- Chap.4 The Software for Life-by-Wire: MOO architectures
 - Efficiency qualities
 - Pay-per-use with prices as resource
 - Configuration optimization
 - MOO software and MOO architectures
- Chap.5 Life-by-Wire

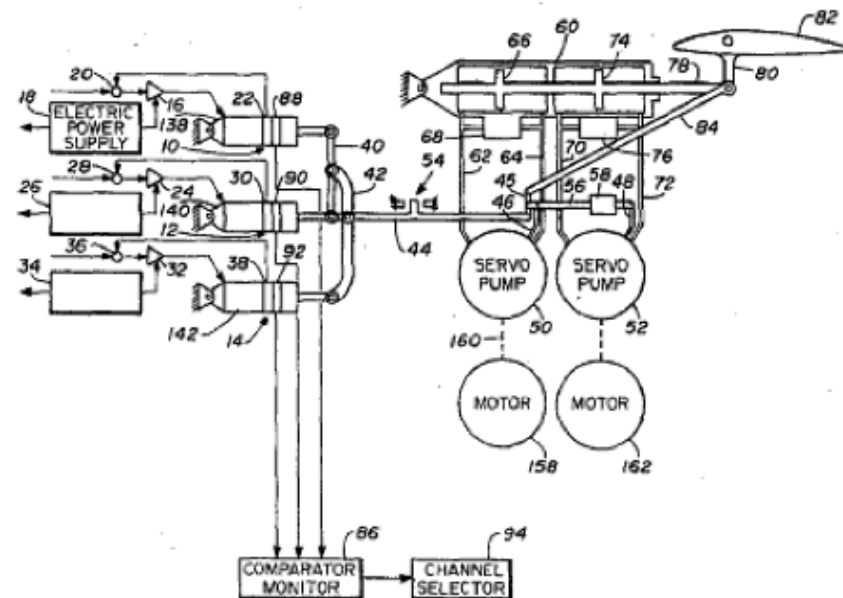


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Chap.1 The Presence: Fly-by-Wire

- **FLY-BY-WIRE**
- William G. Redmond - US Patent 3,679,156, 1972
- <http://www.google.com/patents?hl=de&lr=&vid=USPAT3679156&id=GTswAAAAEBAJ&oi=fnd&dq=fly-by-wire&printsec=abstract#v=onepage&q=fly-by-wire&f=false>



Quality assurance

- Extensive testing

Verification

- Normal mode
- Rescue mode

Certification

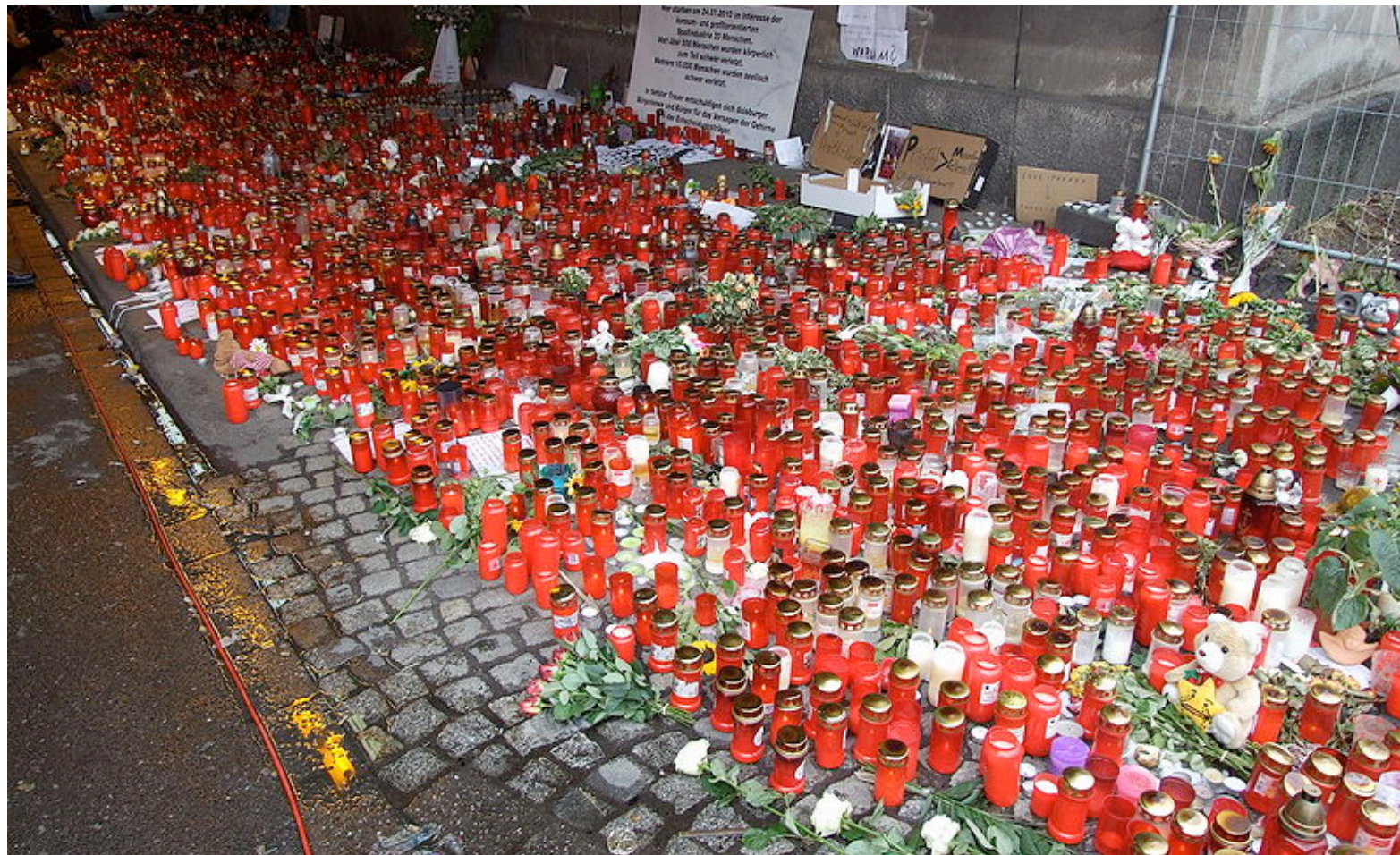
- Certification agencies
- For product liability



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Chap.2 The Near Future: Help-By-Wire



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http://de.wikipedia.org/wiki/Ungl%C3%BCck_bei_der_Loveparade_2010

- „In der Europäischen Union werden jährlich ueber 7.000 Fussgänger und 2.000 Fahrradfahrer bei Verkehrsunfällen getötet. Hunderttausende werden darüber hinaus bei Verkehrsunfällen verletzt.“
- [M. Bischoff. Aktive Sicherheitssysteme fuer den Schutz von Fussgaengern im Strassenverkehr]

Um 4.30 Uhr hatte sich Regina F. auf den Weg zur Arbeit gemacht. Von der Britzer Straße in Schöneweide zum Putzen nach Marienfelde. Sie schaffte es nicht mal bis zum S-Bahnhof.

Viel zu schnell kam ein Kombi um die Ecke gefahren. So schnell, dass das Auto aus der Kurve flog und über den Bürgersteig raste. Durch den Aufprall wurde Regina F. auf die Straße geschleudert. Nur kurz hielt der Mann an, dann gab er Gas und raste davon.



Passive Fail-safe system (FS)

- reach a safe state without power

Quality assurance

- Extensive testing

Verification

- Normal mode
- Safe

Certification

- Certification agencies
- For product liability



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Chap.3 The Mid-Term Future: Move-By-Wire



http://commons.wikimedia.org/wiki/File:Traffic_seen_from_top_of_Arc_de_Triomphe.JPG



Prof. U. Abmann, TU Dresden

Active Fail-safe systems (AFS)

- reach this safe state only with power

Fail-operational systems (FO)

- may fail but continue to work

Quality assurance

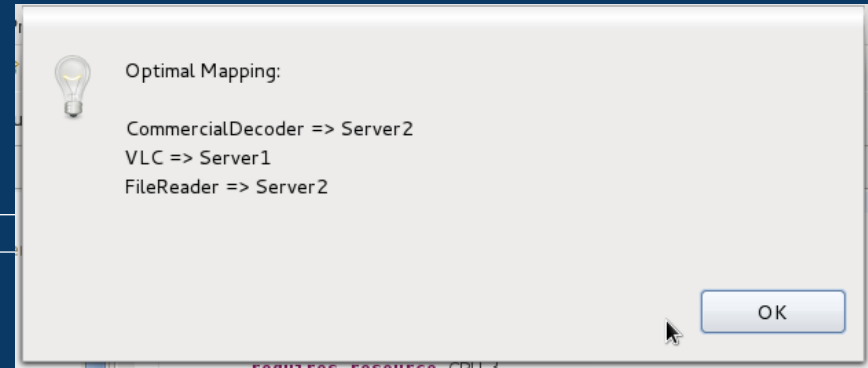
- Extensive testing

Verification

- Normal mode
- Safe

Certification

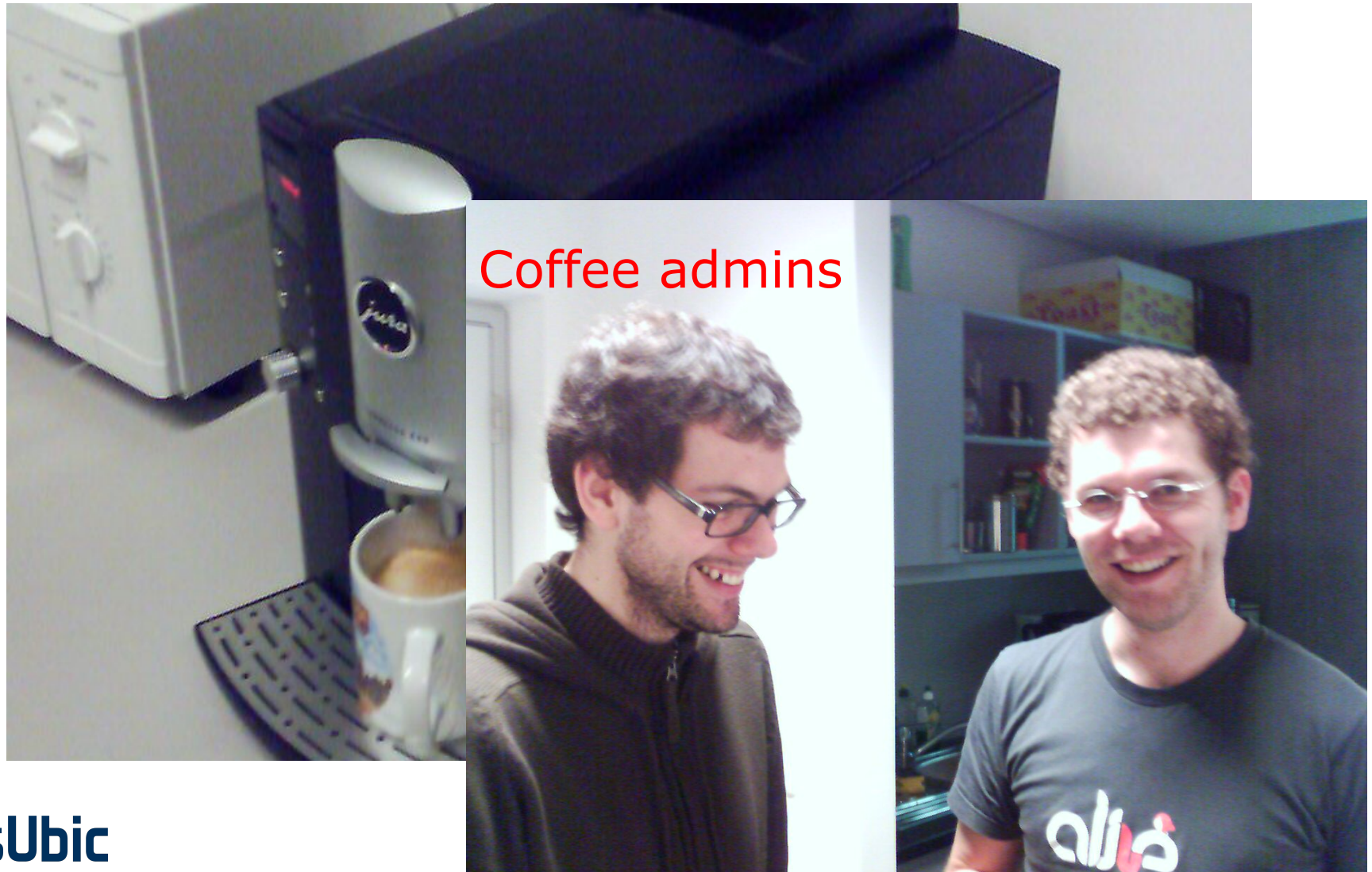
- Certification agencies
- For product liability



Chap.4. The Software for Life-by-Wire: Multi-Objective Optimization Architectures (MOO Architectures)

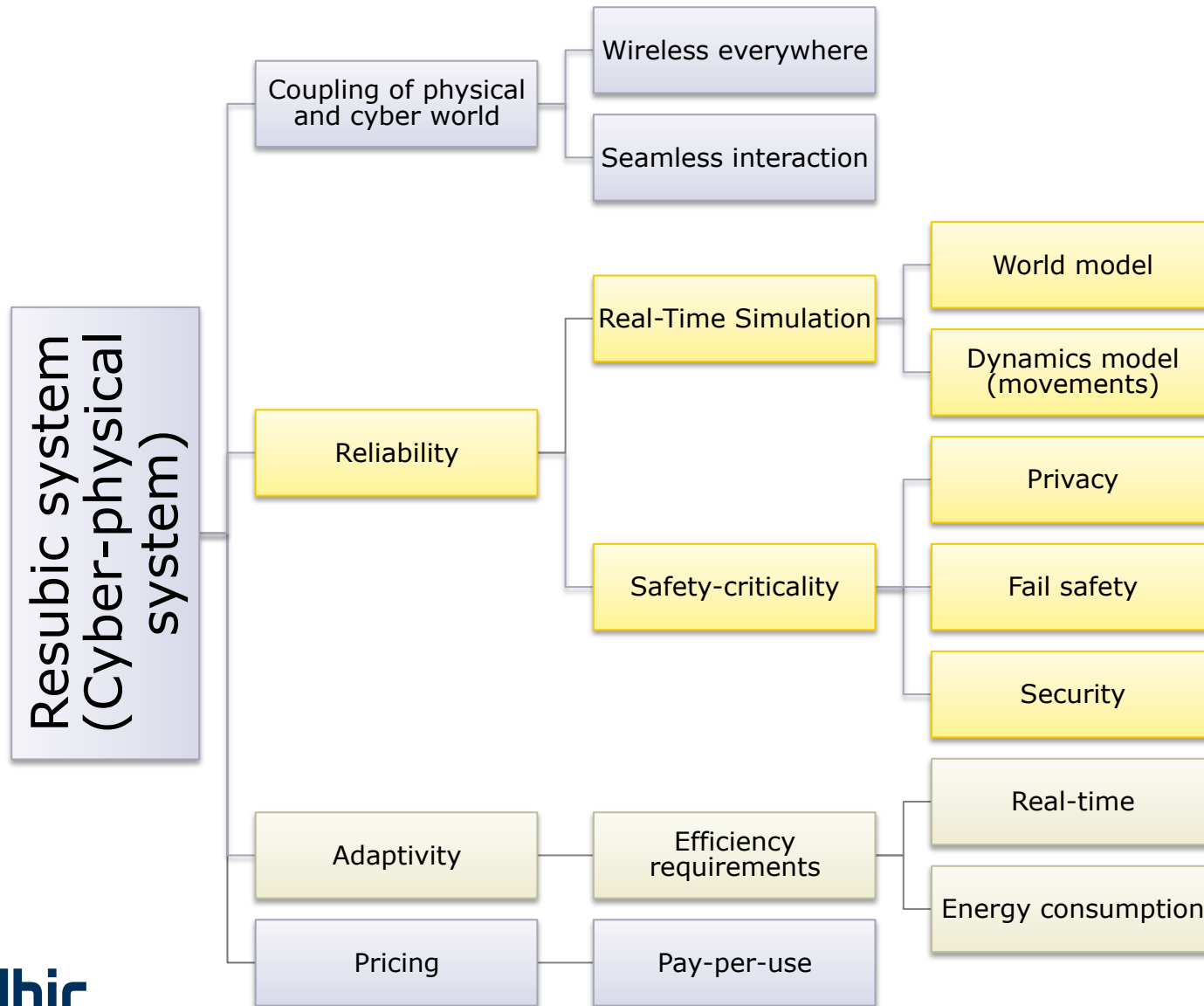
Prof. Uwe Aßmann
Sebastian Götz

Prof. U. Abmann, TU Dresden

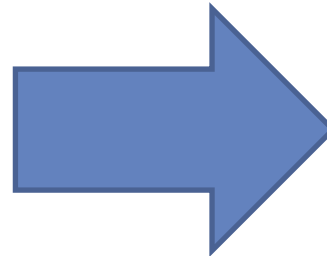




Coffee-by-wire

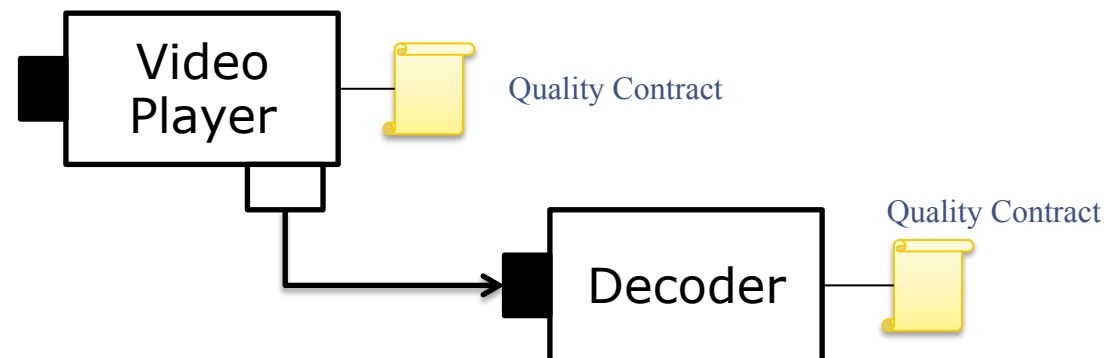


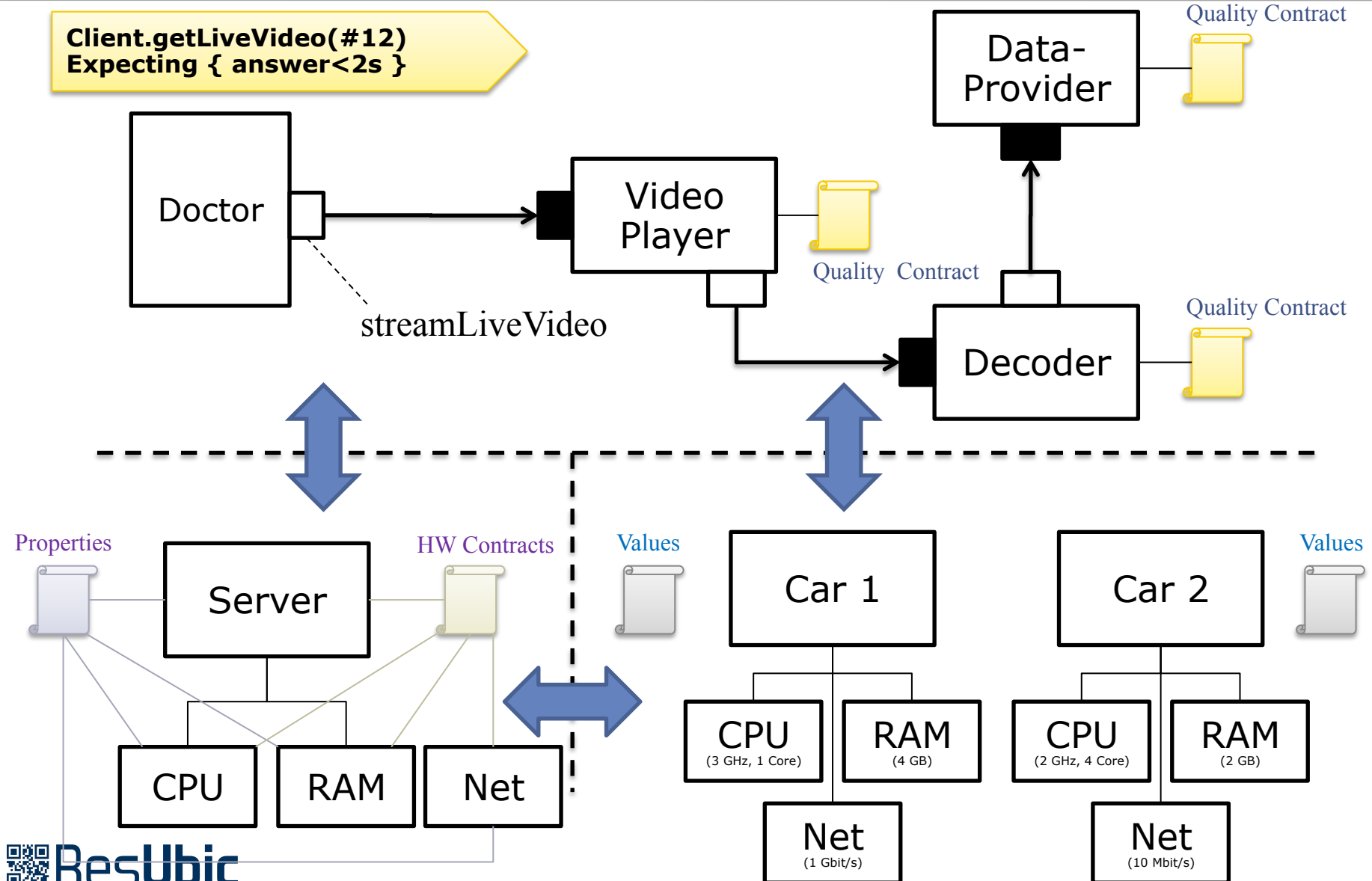
- CPS need adaptability
 - QoS adaptation
 - Context adaptation
- CPS must be reliably composable
 - Functional Contracts
 - Quality Contracts
 - Real-time contracts
 - Energy contracts
 - Safety contracts
 - Spatial dynamics contracts
 - Energy reserve for active fail-safety (AFS)



- Contract Checking of Quality Properties
- Configuration Optimization
 - Determining Optimal Mappings of Software Components onto Hardware Resources
 - Resource minimization (energy)
 - Computation of pricing

- **Temporal path-based quality contracts** are constructs by which the test coverage can be decided statically
- Quality contracts allow for writing statements over time that are verified
 - „in 5ms the system is in a safe state“
 - „the system is life and will never get stuck“
 - „After the door is closed, the train will eventually run“
 - „The robot swarm can still run 25s, so the shutdown to safe state has start in 5s“
 - „This robot cannot be functionally replaced – function of the swarm fails“
- Quality contracts assure that quality features remain during composition of components





- Click to

Multi-Objective Optimized
Software Architectures (MOO)
Cool Component Model (CCM)



Heterogeneous Rich Components
(HRC)
(SPEEDS, CESAR)

contracts

reliability

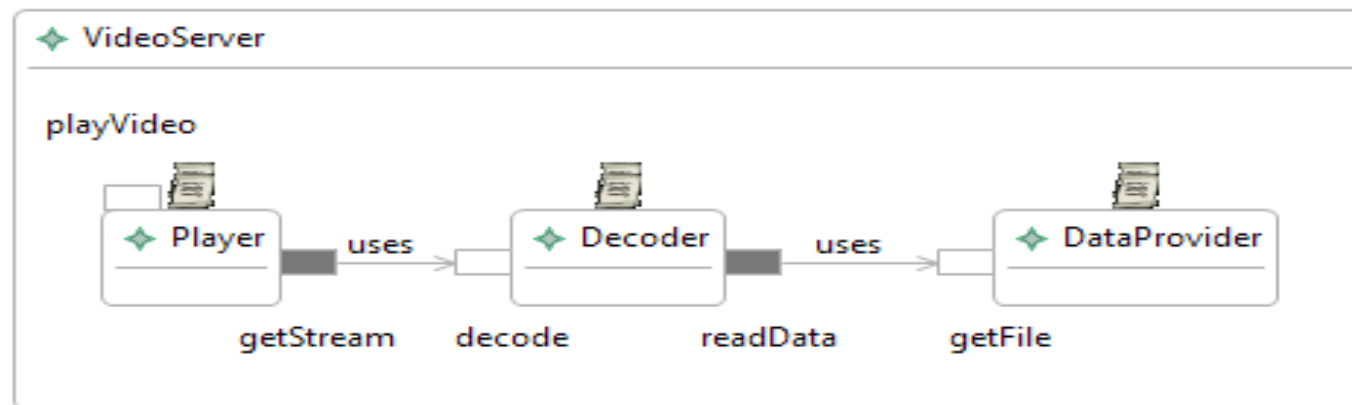
Contract negotiation
(COMQUAD, MADAM, DIVA)

Autonomous Workflow Systems
(OSPP Richly/Aßmann/Lehner)

adaptivity

adaptivity

- Software is described using a CCM Structural Model
 - approx. SysML plus QCL contract language



- SW Components have provided/required ports (e.g. methods of classes)
- Quality-Modes and Implementations of SW Components are described using QCL contracts

```
VideoServer_Player.ecl
1 import ccm [../VideoServer.structure]
2 import ccm [../Server.structure]
3
4 contract VLC implements software Player {
5     mode highQuality {
6         requires component Decoder {
7             bitrate min: 5
8         }
9         requires resource CPU {frequency min: 1500}
10
11         provides framerate min: 20
12     }
13     mode lowQuality {
14         requires component Decoder {
15             bitrate min: 2
16         }
17         requires resource CPU {frequency
18
19         provides framerate min: 10
20     }
21 }
```

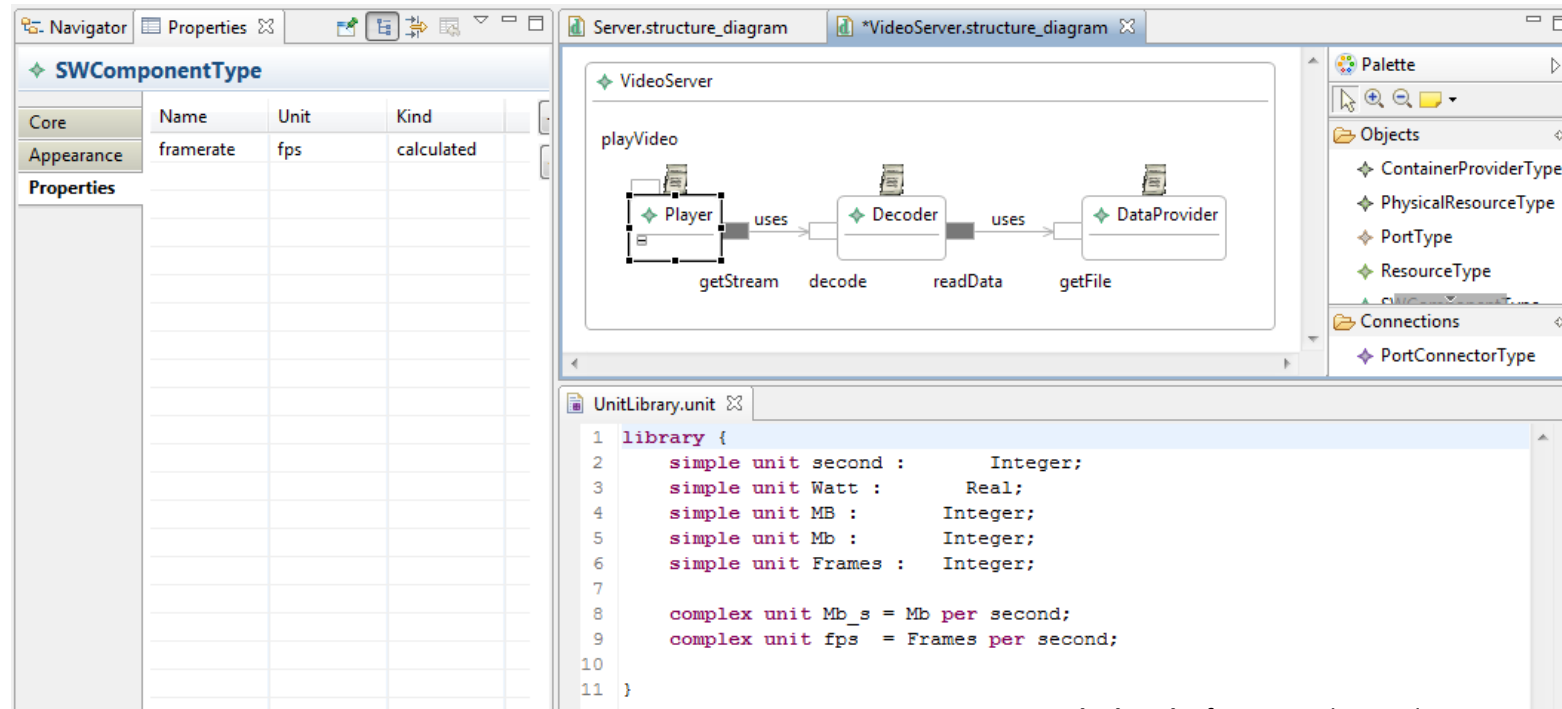
Quality Modes

- SW Dependencies
- HW Dependencies
- Provided Properties

```
VideoServer_Player.ecl  VideoServer_DataProvider.ecl
1 import ccm [../VideoServer.structure]
2 import ccm [../Server.structure]
3
4 contract FileReader implements software DataProvider {
5     mode file {
6         provides throughput min: 20
7         requires resource HDD {
8             throughput min: 20
9         }
10     }
11 }
12
13 contract URLReader implements software DataProvider {
14     mode url {
15         provides throughput min: 5
16         requires resource Network {
17             bandwidth min: 5
18         }
19     }
20 }
```

Where do the properties come from?

Where and how to define HW ?



The screenshot displays a software development environment with three main windows:

- SWComponentType Properties Table:**

Core	Name	Unit	Kind
Appearance	framerate	fps	calculated
- VideoServer Structure Diagram:**

```

graph LR
    subgraph playVideo
        direction LR
        Player[Player] -- uses --> Decoder[Decoder]
        Decoder -- uses --> DataProvider[DataProvider]
    end
    Player -- getStream --> Decoder
    Decoder -- decode --> DataProvider
    DataProvider -- readData --> DataProvider
    DataProvider -- getFile --> DataProvider
    
```
- UnitLibrary.unit Code:**

```

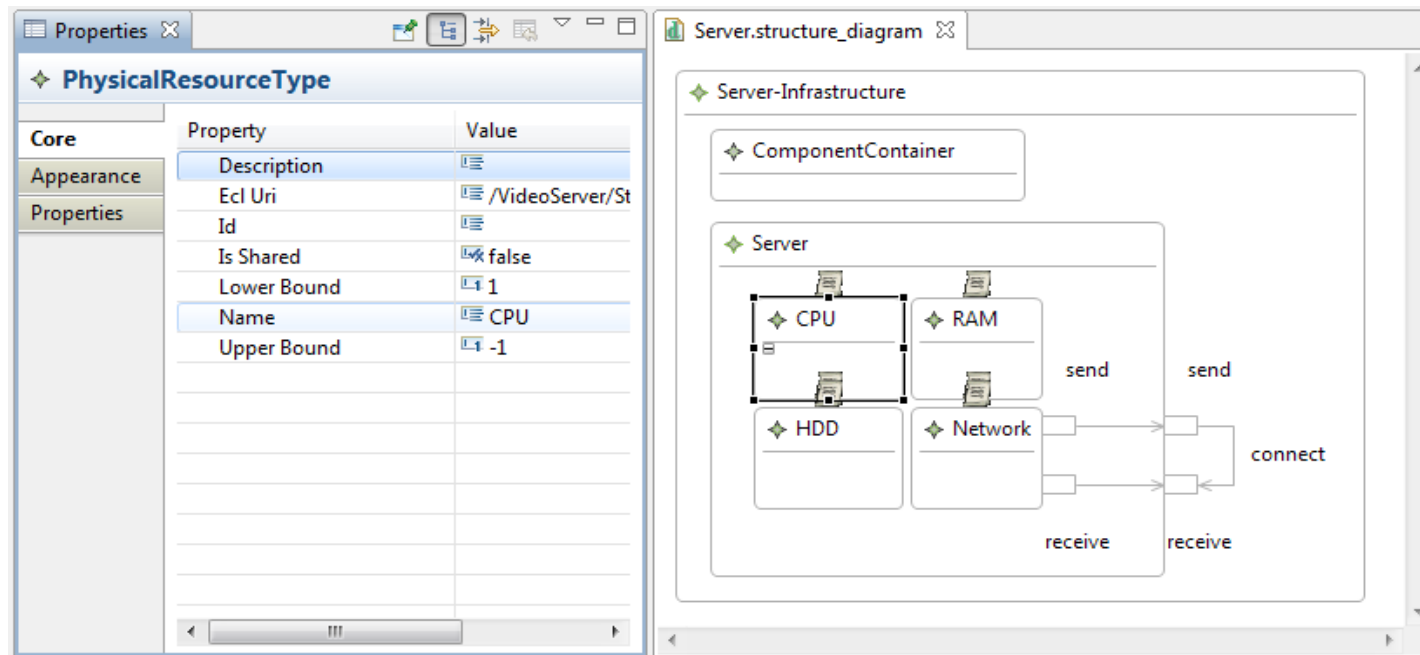
1 library {
2     simple unit second :      Integer;
3     simple unit Watt :       Real;
4     simple unit MB :         Integer;
5     simple unit Mb :         Integer;
6     simple unit Frames :     Integer;
7
8     complex unit Mb_s = Mb per second;
9     complex unit fps = Frames per second;
10
11 }
    
```

calculated: free = total – used
monitored: ResourceMgr.getX()
static: fixed value for instance
 (e.g., size of HDD)

- Properties are defined for each ComponentType.
- Units can be defined
- They are either calculated, monitored or static.

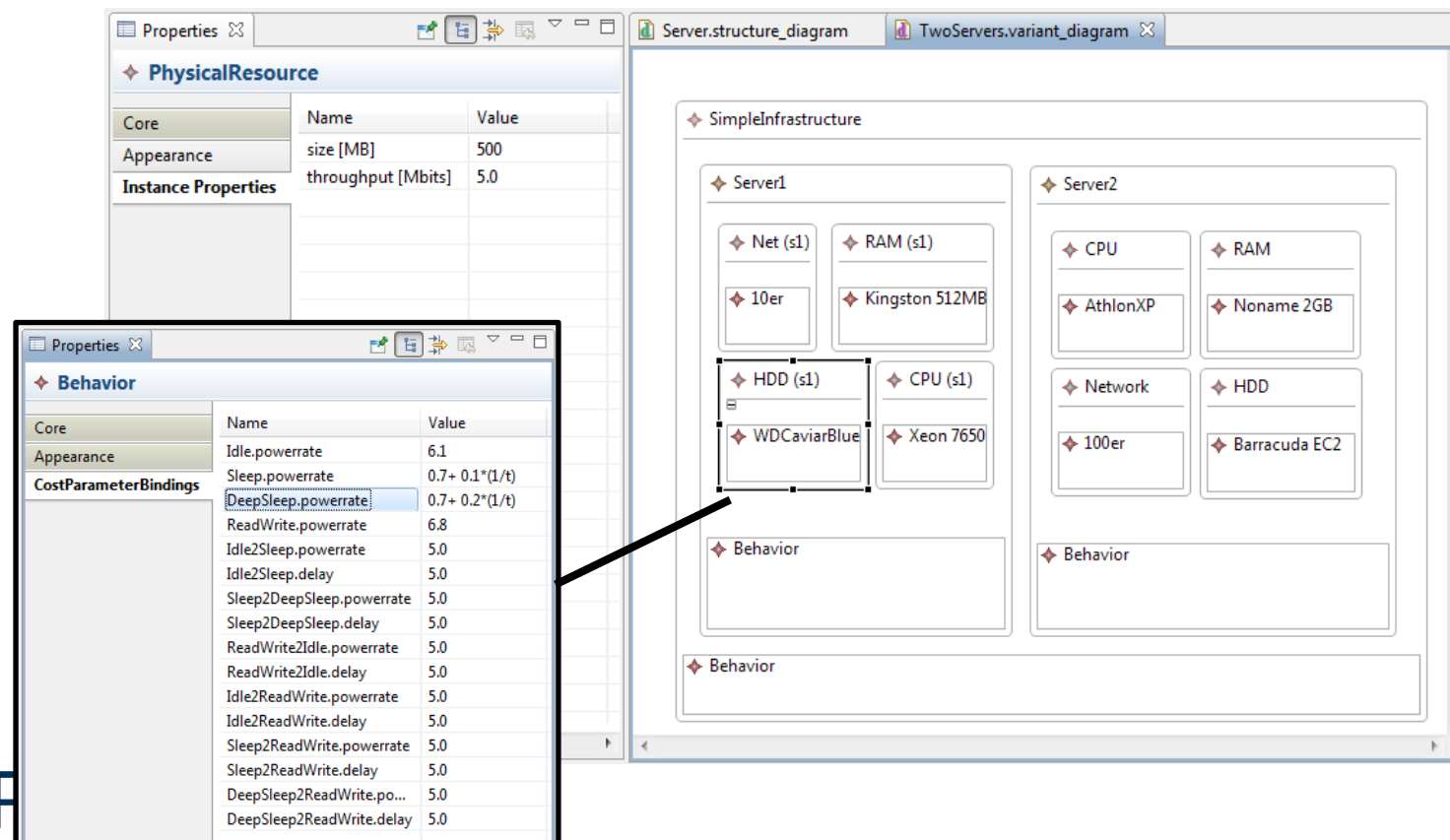
- Hardware Structural Model

- Which ResourceTypes exist
- How they are connected
- Multiplicities - Contracts



- Hardware Variant Model  Concrete Resources (next Slide)
- Hardware Behavior using Energy State Charts

- concrete resources are modeled and connected to structural model
- values of static properties are assigned per resource
- behavior templates are assigned per resource,
- Cost parameters are assigned with values or formulas



The screenshot displays a software interface for hardware variant modeling. It consists of three main parts:

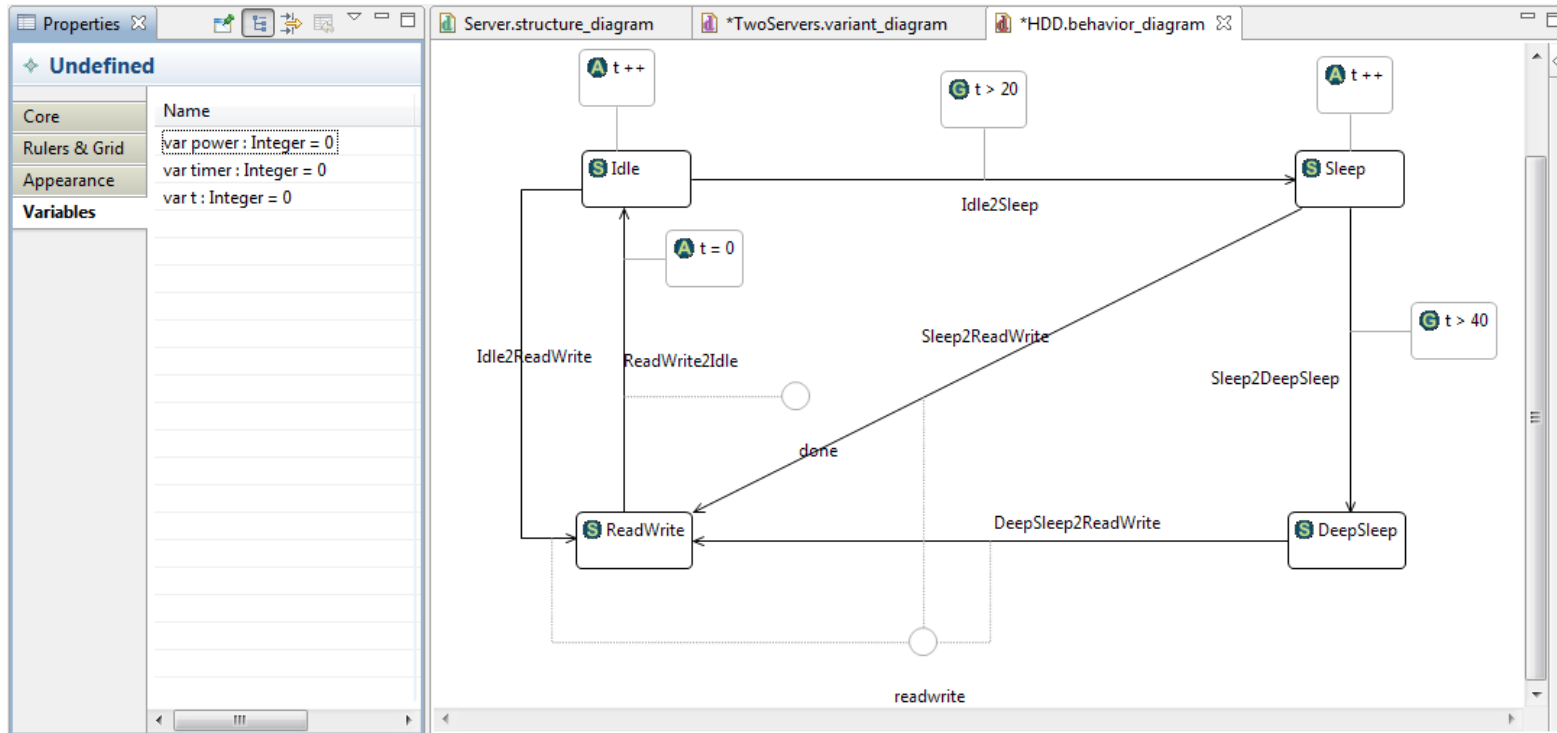
- PhysicalResource Table:** A table showing properties for a physical resource.

Core	Name	Value
Appearance	size [MB]	500
Instance Properties	throughput [Mbits]	5.0
- Behavior Table:** A table showing various power and delay parameters for a behavior.

Core	Name	Value
Appearance	Idle.powerrate	6.1
CostParameterBindings	Sleep.powerrate	$0.7 + 0.1 \cdot (1/t)$
	DeepSleep.powerrate	$0.7 + 0.2 \cdot (1/t)$
	ReadWrite.powerrate	6.8
	Idle2Sleep.powerrate	5.0
	Idle2Sleep.delay	5.0
	Sleep2DeepSleep.powerrate	5.0
	Sleep2DeepSleep.delay	5.0
	ReadWrite2Idle.powerrate	5.0
	ReadWrite2Idle.delay	5.0
	Idle2ReadWrite.powerrate	5.0
	Idle2ReadWrite.delay	5.0
	Sleep2ReadWrite.powerrate	5.0
	Sleep2ReadWrite.delay	5.0
	DeepSleep2ReadWrite.po...	5.0
	DeepSleep2ReadWrite.delay	5.0
- Structural Diagram:** A diagram showing a 'SimpleInfrastructure' containing two servers, 'Server1' and 'Server2'.
 - Server1:** Contains 'Net (s1)' (10er), 'RAM (s1)' (Kingston 512MB), 'HDD (s1)' (WDCaviarBlue), and 'CPU (s1)' (Xeon 7650). It also has a 'Behavior' component.
 - Server2:** Contains 'CPU' (AthlonXP), 'RAM' (Noname 2GB), 'Network' (100er), and 'HDD' (Barracuda EC2). It also has a 'Behavior' component.

An arrow points from the 'DeepSleep.powerrate' entry in the Behavior table to the 'HDD (s1)' component in Server1, indicating that the cost parameter is assigned to that specific resource.





- use cost parameters instead of concrete values
 - reuse for resources of same kind
- can be executed using workload descriptions

Workload Descriptions

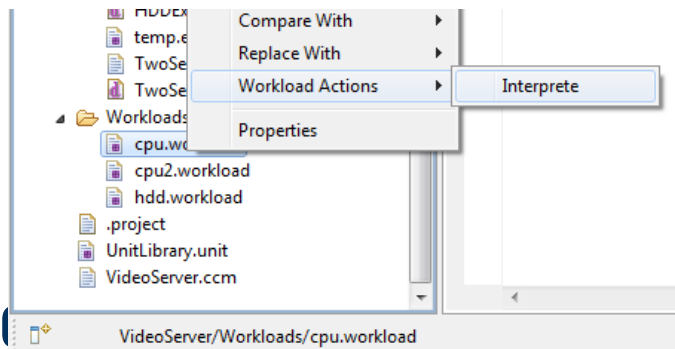
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```

1 import ccm [./VideoServer/Variants/HDDExampleTest.variant]
2
3 Workload TestHDD for WDHDD {
4   readwrite at 2
5   done at 5
6   readwrite at 20
7   done at 30
8   readwrite at 60
9   done at 100
10  readwrite at 150
11  done at 155
12 }
  
```

```

1 import ccm [./VideoServer/Variants/CPUEXampleTest.variant]
2
3 Workload someName for Xeon {
4   work at 5
5   work at 20 for 15
6   work at 50 for 20
7 }
  
```



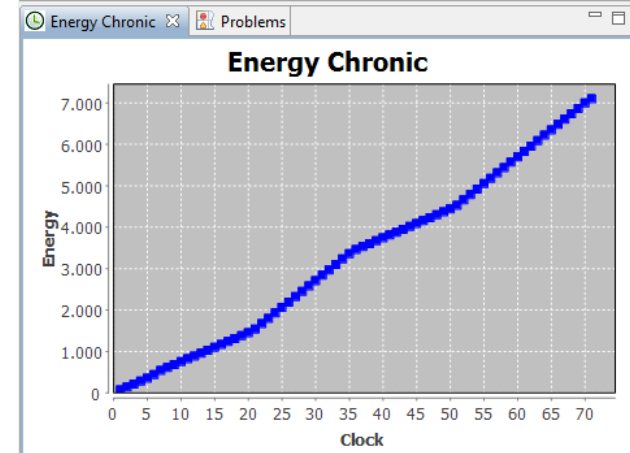
File Explorer: VideoServer/Workloads/cpu.workload

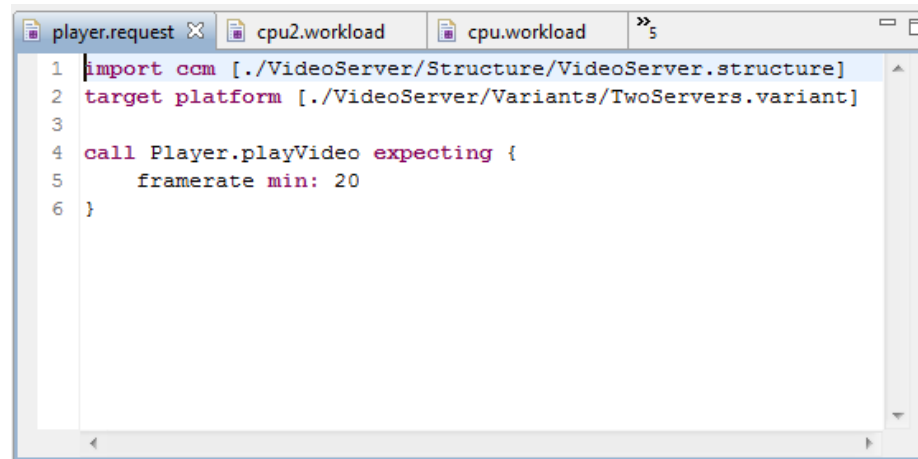
- Compare With
- Replace With
- Workload Actions
 - Interprete
- Properties

Simulation Results

CPU

Clock	Power	State	Config
1	70	Idle	t = 1, x =
2	140	Idle	t = 2, x =
3	210	Idle	t = 3, x =
4	280	Idle	t = 4, x =
5	350	Idle	t = 5, x =
6	440	Busy	t = 5, x = 1
7	540	Idle	t = 1, x = 1
8	610	Idle	t = 2, x = 1
9	680	Idle	t = 3, x = 1
10	750	Idle	t = 4, x = 1
11	820	Idle	t = 5, x = 1
12	890	Idle	t = 6, x = 1
13	960	Idle	t = 7, x = 1
14	1030	Idle	t = 8, x = 1
15	1100	Idle	t = 9, x = 1
16	1170	Idle	t = 10, x = 1
17	1240	Idle	t = 11, x = 1
18	1310	Idle	t = 12, x = 1
19	1380	Idle	t = 13, x = 1
20	1450	Idle	t = 14, x = 1
21	1540	Busy	t = 14, x = 2
22	1670	Busy	t = 14, x = 3
23	1800	Busy	t = 14, x = 4
24	1930	Busy	t = 14, x = 5
25	2060	Busy	t = 14, x = 6
26	2190	Busy	t = 14, x = 7
27	2320	Busy	t = 14, x = 8
28	2450	Busy	t = 14, x = 9
29	2580	Busy	t = 14, x = 10
30	2710	Busy	t = 14, x = 11
31	2840	Busy	t = 14, x = 12





```
player.request x cpu2.workload cpu.workload »5
1 import ccm [./VideoServer/Structure/VideoServer.structure]
2 target platform [./VideoServer/Variants/TwoServers.variant]
3
4 call Player.playVideo expecting {
5   framerate min: 20
6 }
```

- Users call methods on SW Components
- They state their (quality) expectations explicitly
- HW infrastructure is selected

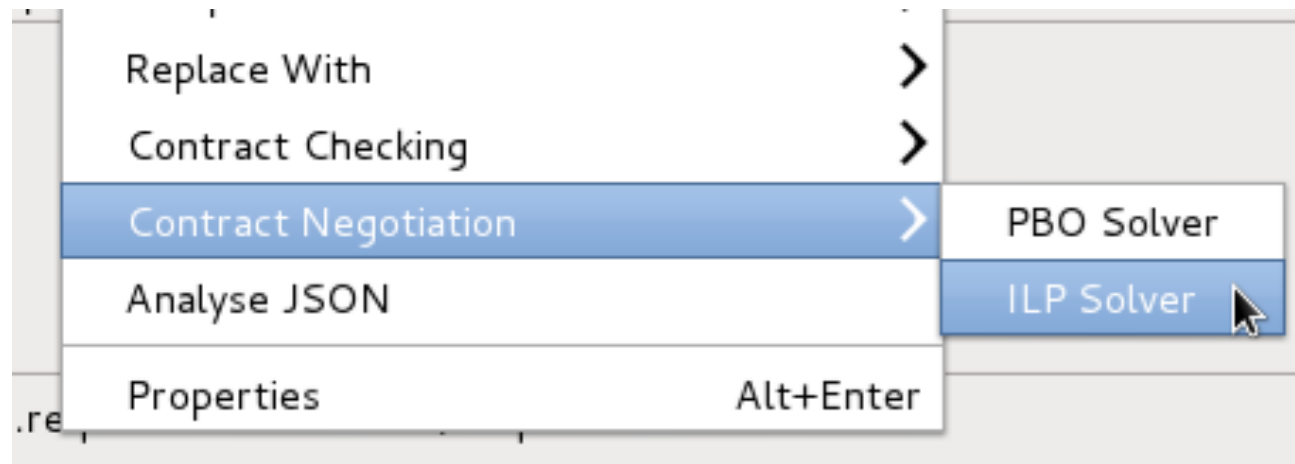
- Then the software configuration is automatically optimized with Multi-Objective-Optimization (MOO)

Question:

- Which implementations of the contracts
- in which quality mode
- on which resources
- give the most efficient configuration?

Multi-Objective Optimization (MOO) with:

- Integer Linear Programming (exact)
- Pseudo-Boolean Optimization (exact)
- Ant Colony Optimization (approx.)
- Simulated Annealing (approx.)
- ...



- An ILP is generated for a certain request on a software component and a hardware variant.

```
VideoServer_Decoder.ecl  player.request.lp  [X]
min: usage#Server1#CPU_[s1]#frequency + usage#Server2#CPU#frequency + usage#Server1#Net_[s1]#bandwidth + usage#Server1#HDD_[s1]#size + usage#Server2#Network#bandwidth
b#VLC#highQuality#Server2 + b#VLC#highQuality#Server1 + b#VLC#lowQuality#Server2 + b#VLC#lowQuality#Server1 = 1;
b#FreeDecoder#fast#Server2 + b#FreeDecoder#fast#Server1 + b#FreeDecoder#slow#Server2 + b#FreeDecoder#slow#Server1 + b#CommercialDecoder#ultrafast#Server2 + b#CommercialDecoder#fast#Server1 + b#CommercialDecoder#slow#Server2 + b#CommercialDecoder#ultrafast#Server1 + b#CommercialDecoder#fast#Server2 + b#CommercialDecoder#slow#Server1 = 1;
b#FileReader#file#Server2 + b#FileReader#file#Server1 + b#URLReader#url#Server2 + b#URLReader#url#Server1 = 1;

usage#Server1#CPU_[s1]#frequency <= 2000.0;
usage#Server1#CPU_[s1]#frequency >= 0;
usage#Server1#CPU_[s1]#frequency = 1500.0 b#VLC#highQuality#Server1 + 1500.0 b#FreeDecoder#fast#Server1 + 1000.0 b#FreeDecoder#slow#Server1 + 800.0 b#CommercialDecoder#ultrafast#Server1 + 800.0 b#CommercialDecoder#fast#Server1 + 800.0 b#CommercialDecoder#slow#Server1;
usage#Server1#RAM_[s1]#size <= 512.0;
usage#Server1#RAM_[s1]#size >= 0;
usage#Server1#RAM_[s1]#size = 512.0 b#FreeDecoder#fast#Server1 + 256.0 b#FreeDecoder#slow#Server1 + 128.0 b#CommercialDecoder#slow#Server1 + 512.0 b#CommercialDecoder#ultrafast#Server1 + 512.0 b#CommercialDecoder#fast#Server1 + 512.0 b#CommercialDecoder#slow#Server1;
usage#Server1#Net_[s1]#bandwidth <= 10.0;
usage#Server1#Net_[s1]#bandwidth >= 0;
usage#Server1#Net_[s1]#bandwidth = 5.0 b#URLReader#url#Server1;
usage#Server1#HDD_[s1]#size <= 500.0;
usage#Server1#HDD_[s1]#size >= 0;
usage#Server1#HDD_[s1]#throughput_w <= 5.0;
usage#Server1#HDD_[s1]#throughput_w >= 0;
usage#Server2#CPU#frequency <= 1000.0;
usage#Server2#CPU#frequency >= 0;
usage#Server2#CPU#frequency = 1500.0 b#VLC#highQuality#Server2 + 1500.0 b#FreeDecoder#fast#Server2 + 1000.0 b#FreeDecoder#slow#Server2 + 800.0 b#CommercialDecoder#ultrafast#Server2 + 800.0 b#CommercialDecoder#fast#Server2 + 800.0 b#CommercialDecoder#slow#Server2;
usage#Server2#RAM#size <= 2048.0;
usage#Server2#RAM#size >= 0;
usage#Server2#RAM#size = 512.0 b#FreeDecoder#fast#Server2 + 256.0 b#FreeDecoder#slow#Server2 + 128.0 b#CommercialDecoder#slow#Server2 + 512.0 b#CommercialDecoder#ultrafast#Server2 + 512.0 b#CommercialDecoder#fast#Server2 + 512.0 b#CommercialDecoder#slow#Server2;
usage#Server2#Network#bandwidth <= 100.0;
usage#Server2#Network#bandwidth >= 0;
usage#Server2#Network#bandwidth = 5.0 b#URLReader#url#Server2;
usage#Server2#HDD#size <= 150.0;
usage#Server2#HDD#size >= 0;
usage#Server2#HDD#throughput_w <= 25.0;
usage#Server2#HDD#throughput_w >= 0;
bitrate = 10.0 b#CommercialDecoder#fast#Server1 + 8.0 b#FreeDecoder#fast#Server2 + 3.0 b#FreeDecoder#slow#Server2 + 10.0 b#CommercialDecoder#fast#Server2 + 3.0 b#CommercialDecoder#slow#Server2 + 10.0 b#CommercialDecoder#ultrafast#Server1 + 10.0 b#CommercialDecoder#fast#Server2 + 10.0 b#CommercialDecoder#slow#Server2 + 10.0 b#CommercialDecoder#ultrafast#Server2 + 10.0 b#CommercialDecoder#fast#Server1;
throughput = 5.0 b#URLReader#url#Server1 + 20.0 b#FileReader#file#Server2 + 5.0 b#URLReader#url#Server2 + 20.0 b#FileReader#file#Server1;
framerate = 10.0 b#VLC#lowQuality#Server1 + 10.0 b#VLC#lowQuality#Server2 + 20.0 b#VLC#highQuality#Server2 + 20.0 b#VLC#highQuality#Server1;
bitrate >= 2.0 b#VLC#lowQuality#Server1 + 2.0 b#VLC#lowQuality#Server2 + 5.0 b#VLC#highQuality#Server2 + 5.0 b#VLC#highQuality#Server1;
throughput >= 5.0 b#CommercialDecoder#fast#Server1 + 5.0 b#FreeDecoder#fast#Server2 + 2.0 b#FreeDecoder#slow#Server2 + 5.0 b#CommercialDecoder#fast#Server2 + 2.0 b#CommercialDecoder#slow#Server2 + 5.0 b#CommercialDecoder#ultrafast#Server1 + 5.0 b#CommercialDecoder#fast#Server2 + 5.0 b#CommercialDecoder#slow#Server2 + 5.0 b#CommercialDecoder#ultrafast#Server2 + 5.0 b#CommercialDecoder#fast#Server1;
framerate >= 20.0;

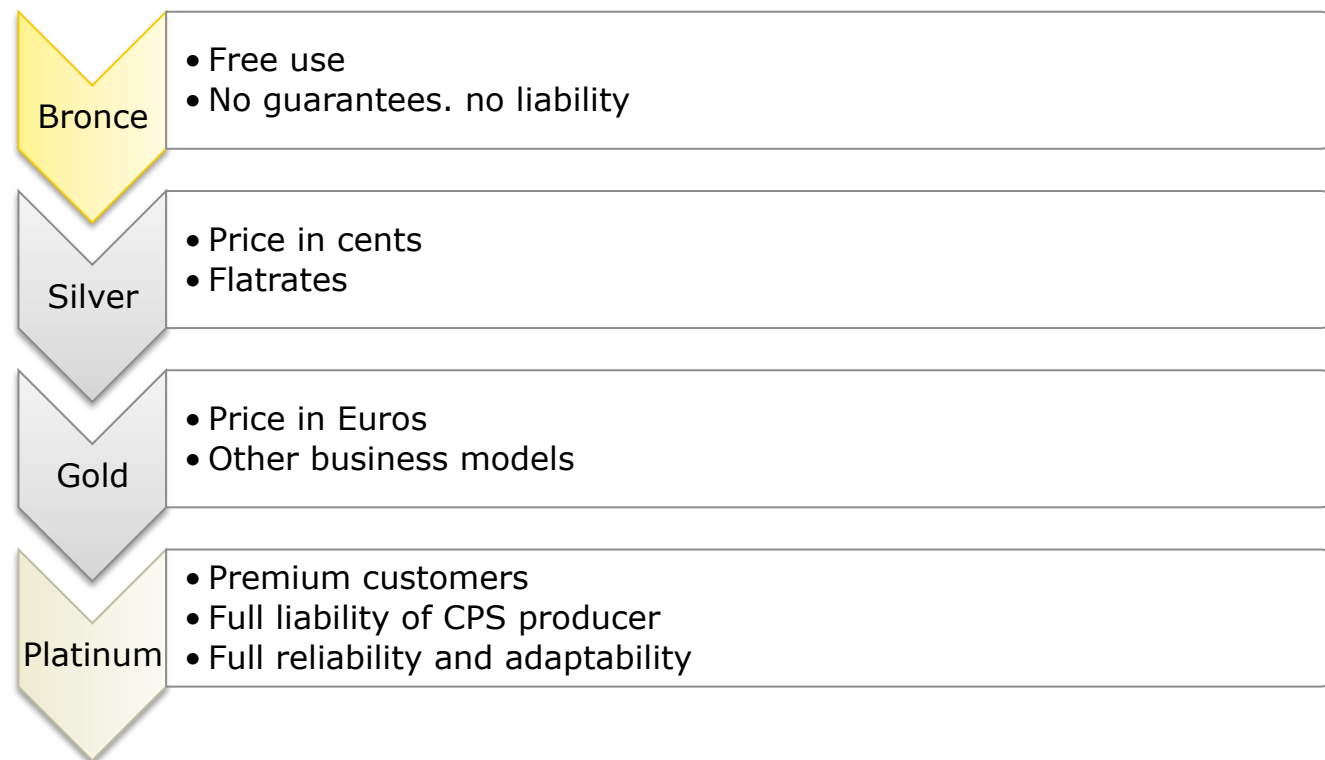
binary b#CommercialDecoder#fast#Server1, b#CommercialDecoder#fast#Server2, b#URLReader#url#Server1, b#FileReader#file#Server1, b#FileReader#file#Server2, b#URLReader#url#Server2, b#VLC#lowQuality#Server1, b#VLC#lowQuality#Server2, b#VLC#highQuality#Server1, b#VLC#highQuality#Server2, b#FreeDecoder#fast#Server1, b#FreeDecoder#fast#Server2, b#FreeDecoder#slow#Server1, b#FreeDecoder#slow#Server2, b#CommercialDecoder#ultrafast#Server1, b#CommercialDecoder#ultrafast#Server2, b#CommercialDecoder#fast#Server1, b#CommercialDecoder#fast#Server2, b#CommercialDecoder#slow#Server1, b#CommercialDecoder#slow#Server2, b#URLReader#url#Server1, b#URLReader#url#Server2, b#FileReader#file#Server1, b#FileReader#file#Server2
```

Optimal Mapping:

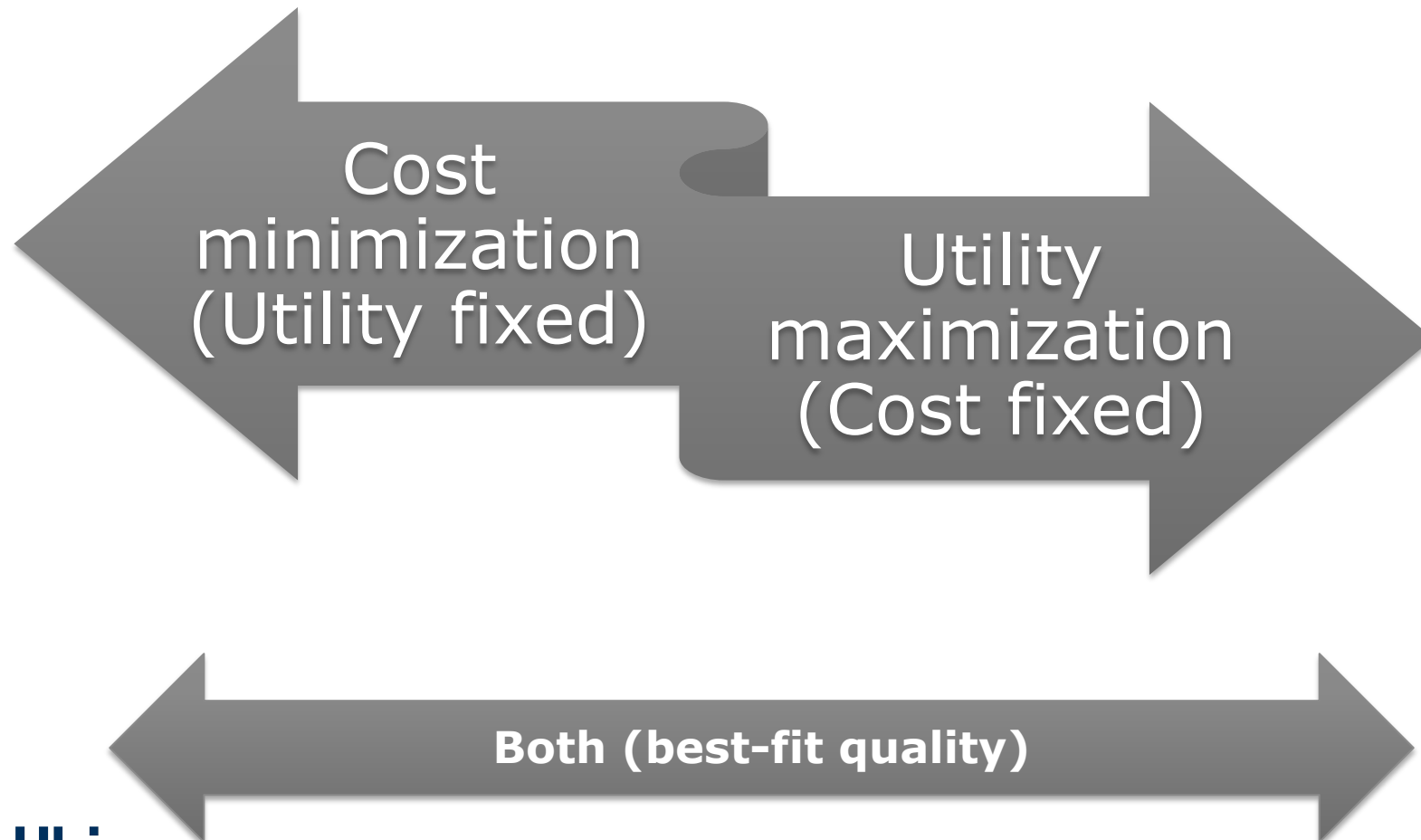
- CommercialDecoder => Server2
- VLC => Server1
- FileReader => Server2

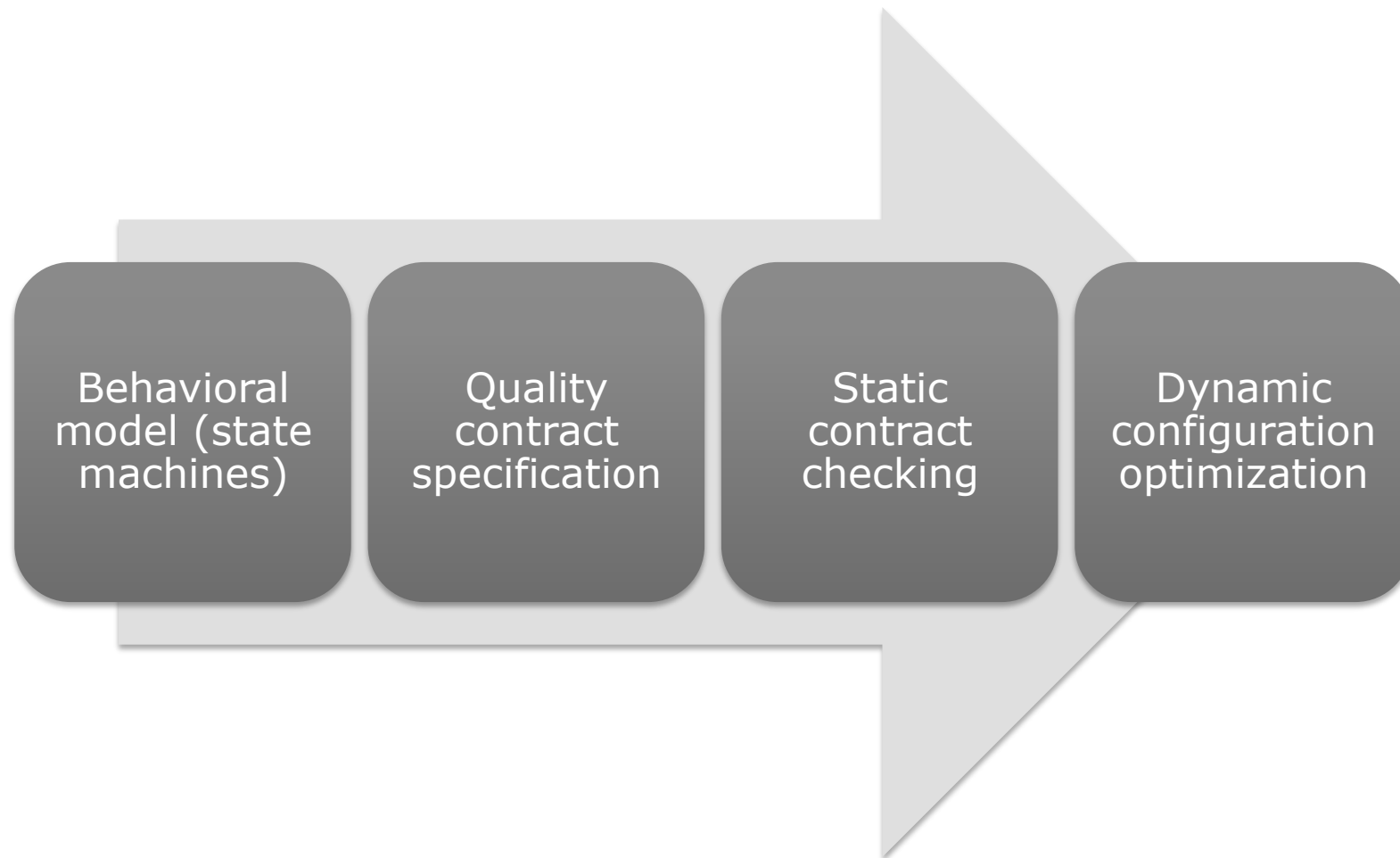
OK

- Prices are costs
- Utility per price is an efficiency feature
- Prices can be negotiated automatically
- Using a CPS will be possible with micro-pay-per-use



- All efficiency requirements can be handled





4.1 Seamless Interaction with CPS

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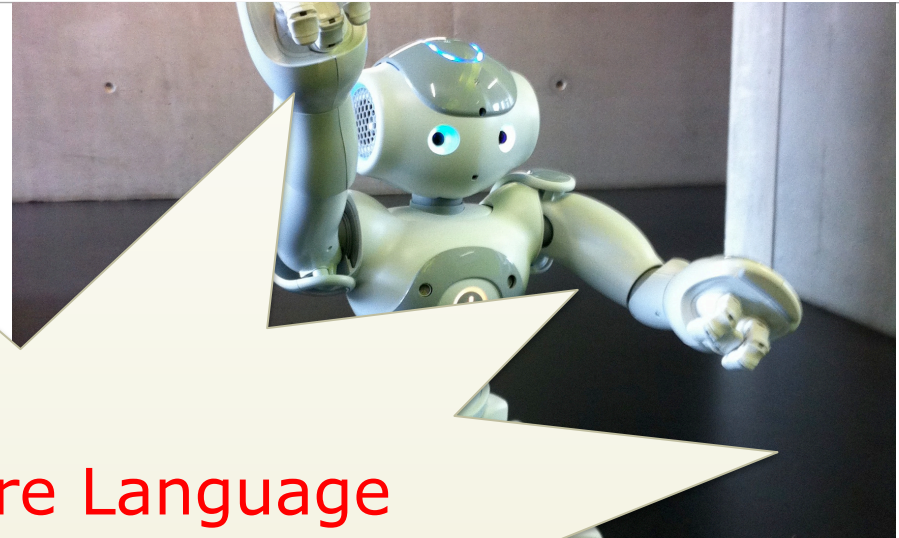
But, how do I teach Chuck to bring me a coffee?

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How do I teach Chuck to bring me a coffee?

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with a Gesture Language



Gestures in Cars

Gestures in Public

Gestures in Factories

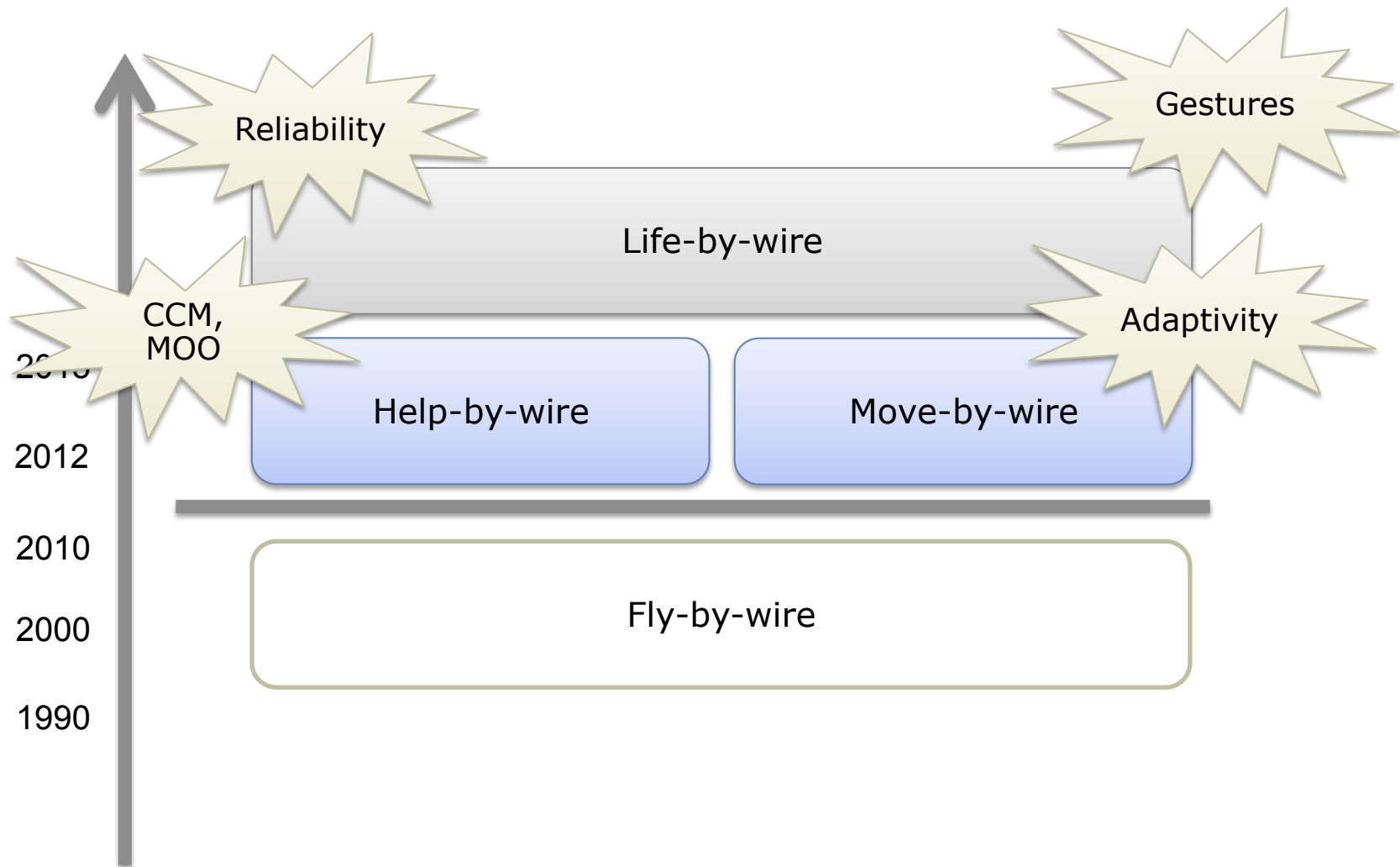
Context modeling

User editable

Semantics



Chap.5. Life-by-Wire



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- Andreas.richter@tu-dresden.de
- Wolfgang.lehner@tu-dresden.de
- Sebastian.goetz@tu-dresden.de

- [Acatech] Agenda CPS. Acatech Studie März 2012.